

THE (NOT SO PROFITABLE) PATH TOWARDS AUTOMATED HEAP EXPLOITATION

Thaís Moreira Hamasaki

 [barbieAuglend](#)

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DISCLAIMER

*This research was accomplished by
me in my personal capacity during
my spare time.*

DON'T BE TOO JUDGEMENTAL PLEASE! :)



full disclosure: I am NOT a vulnerability researcher!



ABOUT ME

```
echo 'Stare at binaries during the day';  
  
echo 'Blackhoodie - Core Organizer and Board Member';  
echo 'HackLu`s program`s committee';  
echo 'Disobey`s Lead of Technical Content';  
echo 'x86 Assembly & RE101 - Lead of both groups @chaosdorf';  
echo 'Logical Programming, RE, static analysis, Mountaineering FTW';  
echo 'Wannabe "Karaoke" singer';  
  
echo 'Stare at binaries by night';
```

BlackHoodie



WHAT AM I GOING TO TALK ABOUT?

- **constraint logic programming (CLP)**
- **solvers**
- **static analysis scalability**
- **the memory**
- **oh yeah, heaps...**



S..A..T.. WHAT?



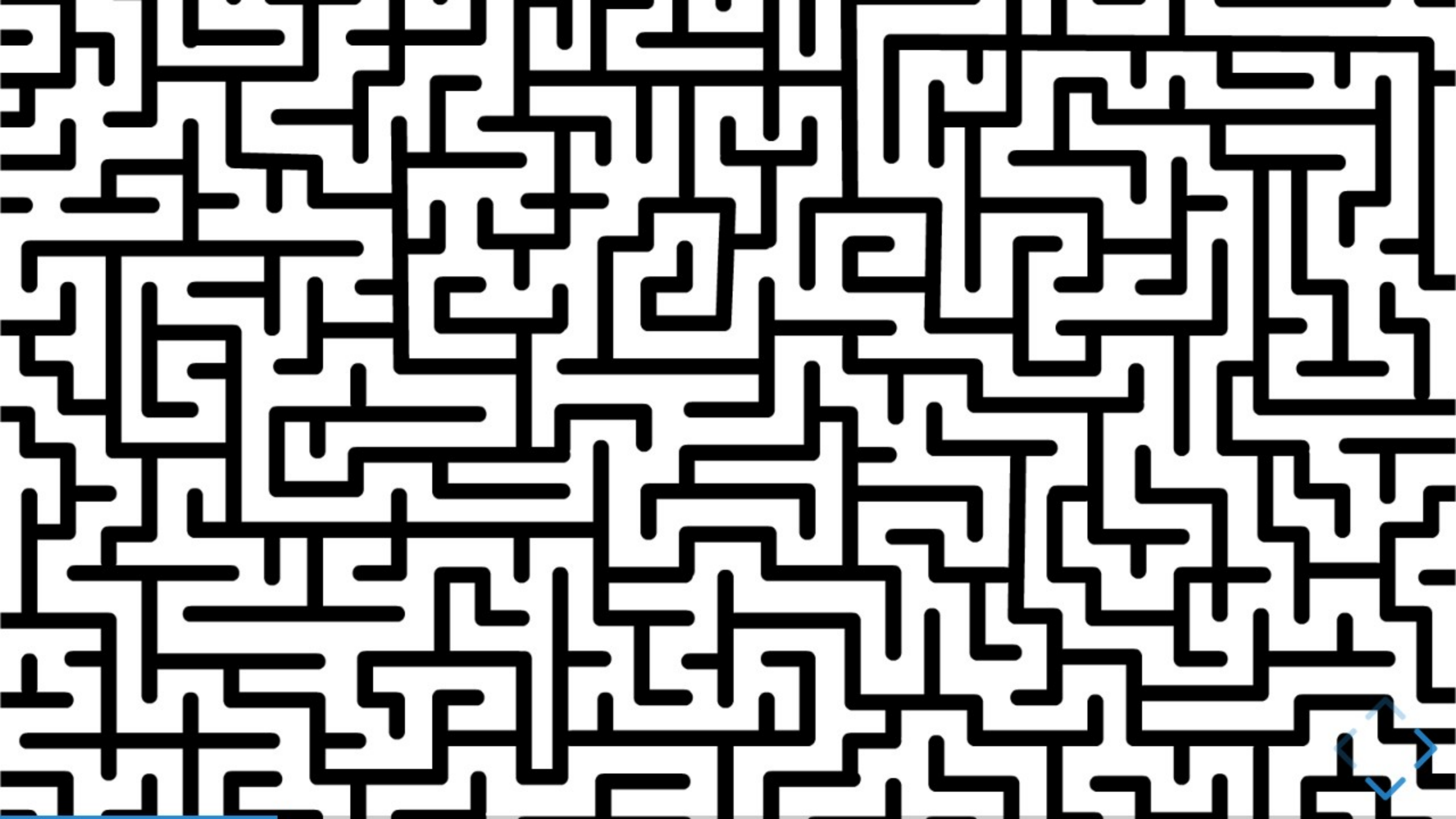
- **Solver!**
- **Satisfiability Modulo Theories (SMT)**





CONSTRAINTS

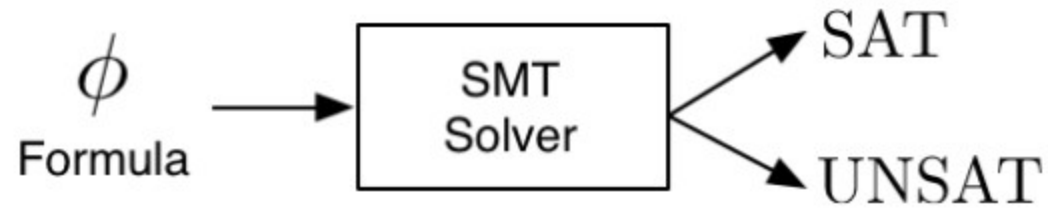




"Constraint programming represents one of the closest approaches computer science has yet made to the Holy Grail of programming: the user states the problem, the computer solves it." Eugene C. Freuder, Constraints, April 1997



AUTOMATED THEOREM PROVING



- **Hardware and Software → Large-scale verification**
- **Languages specification and Computing proof obligations**



SYMBOLIC EXECUTION



IT LOOKS LIKE THAT ...

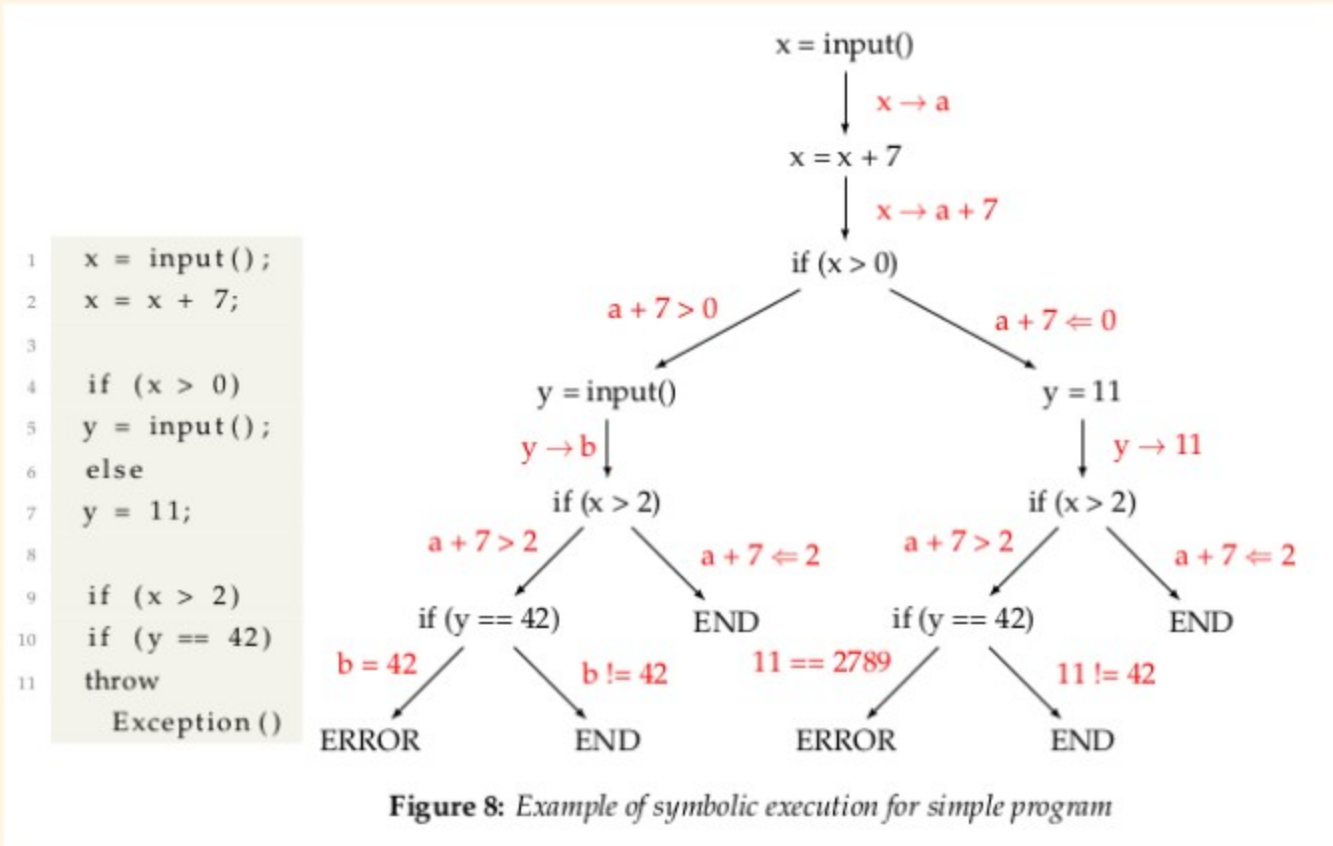


Figure 8: Example of symbolic execution for simple program

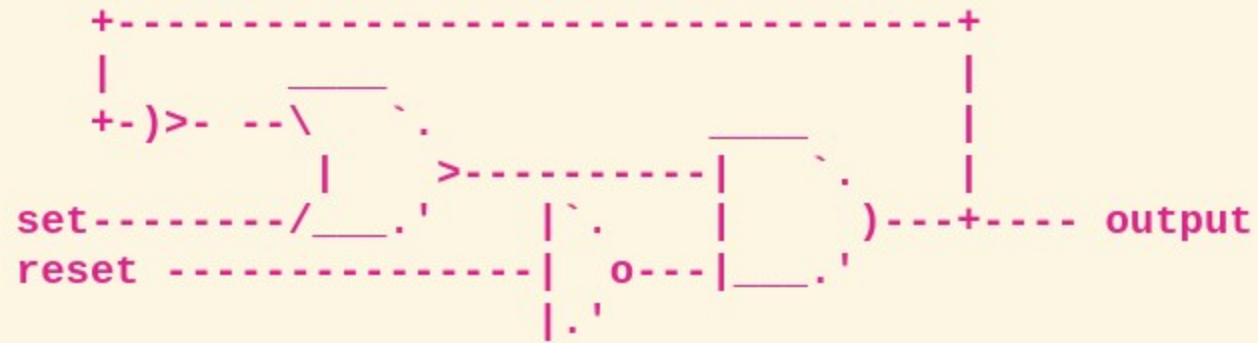


HOW IT WORKS

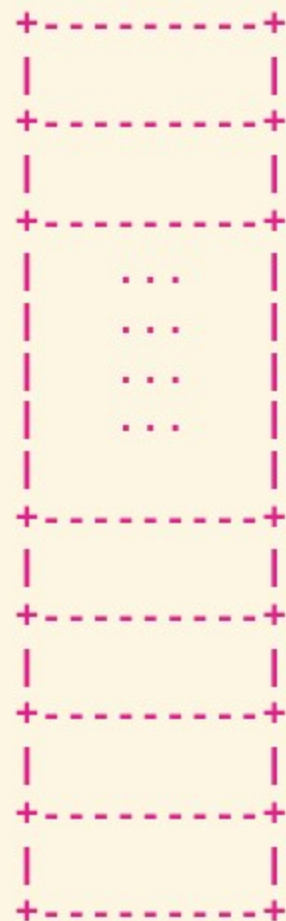
- **Create a process ($pc = 0$, $state = []$)**
- **Add the process (pc , $state$) to the domain system D**
- **while D not empty:**
 - **Remove process (pc , $state$) from system**
 - **Execute it until the next branching point**
 - **If both paths are feasible, add both to D**
 - **if just one is feasible, add the feasible path and the negation of the not feasible path to D**



THE LOGIC GATES OF THE MEMORY



50 SHADES OF MEMORY



STACK

the IMPLEMENTATION

HEAP



A HEAP OF INFORMATION



APPLICATIONS



MALWARE ANALYSIS



- **Obfuscation**
- **Compiler optimizations**
- **Crypto-analysis**



BUG HUNTING



- **Fuzzing**
- **Code verification**
- **Binary Analysis**



EXPLOITATION



- **PoC (Proof of Concept)**
- **AEF (Automated Exploit Framework)**
- **APG (Automated Payload Generation)**



WHAT WE ARE LOOKING FOR



- **Vulnerable**



WHAT WE ARE LOOKING FOR



- **Vulnerable *AND* Exploitable**



HOW TO CRASH !



AUTOMATION OUT THERE



- **Exploratory testing**
- **Dynamic taint analysis**
- **Abstract interpretation**



AUTOMATION OUT THERE



- **Klee**
 - **Open source symbolic executor**
 - **Runs on top of LLVM**
- **Manticore**
 - **Symbolic execution**
 - **Taint analysis**
 - **Binary instrumentation**



TOOL OF CHOICE



FORWARD

SYMBOLIC

EXECUTION



EXPLOIT GENERATION



Find a bug — easy right?

Def: Vulnerable Path for input ϵ is
 $\Pi_{\text{(vulnerability)}}(\epsilon)$



PLAN



**Theorem: Given a program,
automatically find vulnerabilities
and generate exploits for them.**



PLAN



**Theorem: Given a program,
automatically find vulnerabilities
and generate exploits for them.**

- **direct influence**



PLAN



**Theorem: Given a program,
automatically find vulnerabilities
and generate exploits for them.**

- **indirect influence**

```
malloc(strlen(user_input));
```



EXPLOIT GENERATION



Check if it is exploitable

Not that easy anymore...

$\Pi(\text{vulnerability})(\epsilon) \wedge \Pi(\text{exploit})(\epsilon) = \text{true}$

Where $\Pi(\text{exploit})(\epsilon)$ is the attacker's logic



EXPLOIT GENERATION



Implement $\Pi_{\text{exploit}}(\epsilon)$

for a really special case $\Pi_{\text{vulnerability}}(\epsilon) \wedge$
 $\Pi_{\text{exploit}}(\epsilon)$

and then it works MOST of the times

is it really automated then?



EXPLOIT GENERATION



EXTRA: Evaluate

- Find the HEAP
- Exploit Verification
- State Space Explosion
- Environment Definition



LIMITATIONS



THEORETICAL #1

Rice's Theorem

Theorem

Let L be a subset of strings representing Turing machines, where

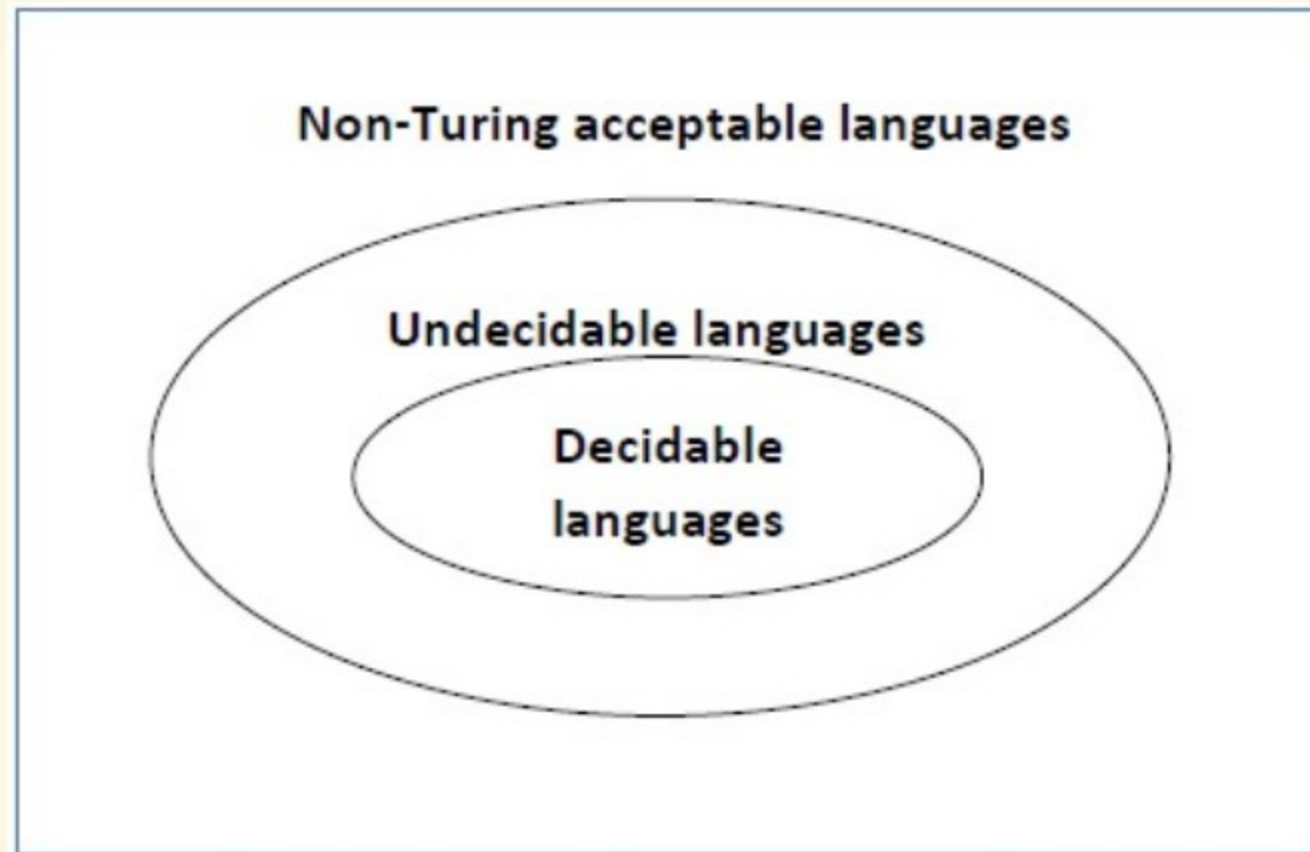
1. If M_1 and M_2 recognize the same language, then either $\langle M_1 \rangle, \langle M_2 \rangle \in L$ or $\langle M_1 \rangle, \langle M_2 \rangle \notin L$.

2. $\exists M_1, M_2$ s.t. $\langle M_1 \rangle \in L$ and $\langle M_2 \rangle \notin L$.

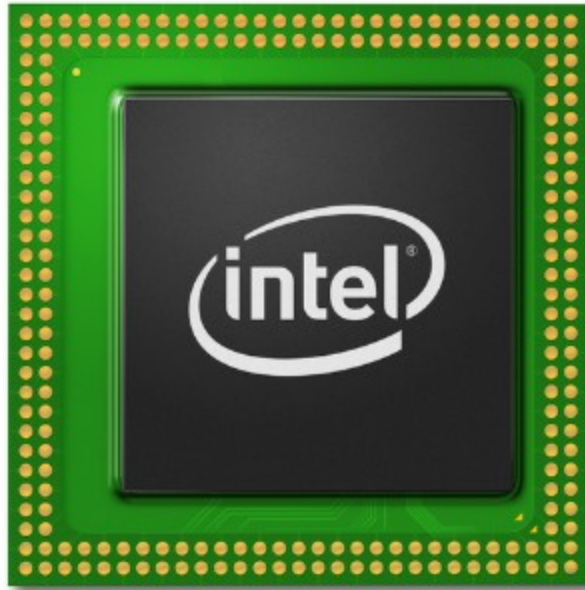
Then L is undecidable.



THEORETICAL #2



PRACTICAL



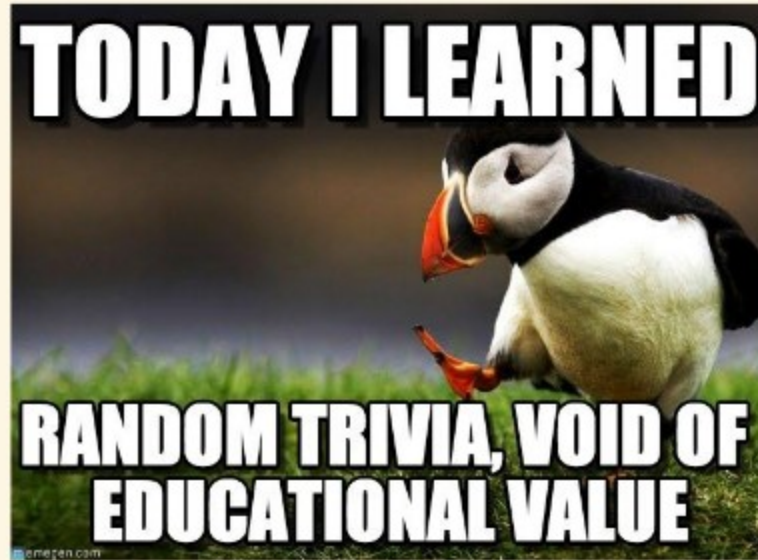
Remember...



CONCLUSION



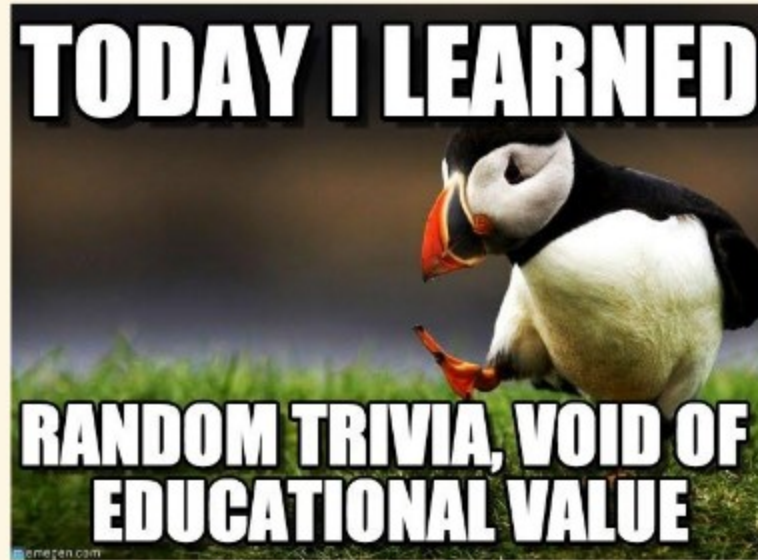
LEARNINGS / TAKE AWAY



- Symbolic execution is a powerful tool ~~while~~ analysing ~~malware~~ for vulnerability research
- SMT solvers can reason and generate exploits



LEARNINGS / TAKE AWAY



- Symbolic execution can be ~~is~~ a powerful tool for vulnerability research
- SMT solvers can reason and generate exploits



WORK DONE:



- a binary garbage-code eliminator, a XOR search, some "cryptographic" algorithm breaker, a generic unpacker, a binary structure recognizer, a C++ class hierarchy reconstructor.





WORKING ON ...



- **specialized constraint inference assistant for computer security problems**



ACKNOWLEDGMENTS

-  Sean Heelan
 - Automated Heap Layout Manipulation for Exploitation (Heelan et al. to appear in Usenix Security 2018)
 - and his time;
 - and inspiration!
-  Marion Marschalek
- Heap Models for Exploit Systems (Vanegue, Langsec 2015)
- and the Intel Documentation I think ... ?



QUESTIONS?



 @barbieAuglend

 barbie@barbieauglend.re

